

A Dynamically Reconfigurable Processor Architecture

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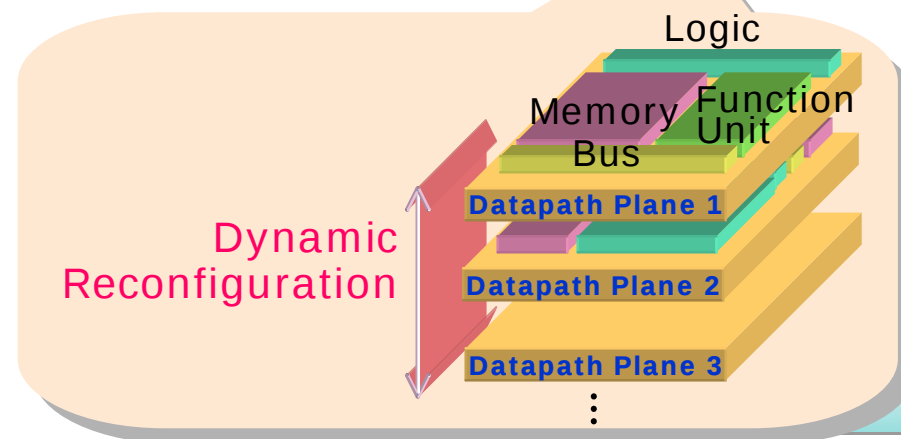
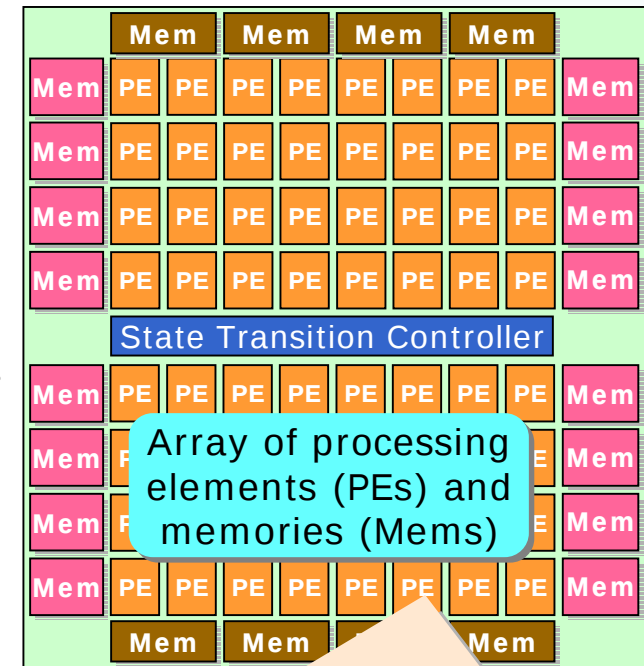
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- ❖ Introduction
- ❖ DRP architecture
- ❖ DRP compiler
- ❖ DRP Application development system
- ❖ Sample applications
- ❖ Future roadmap and summary

Introduction: What is DRP?

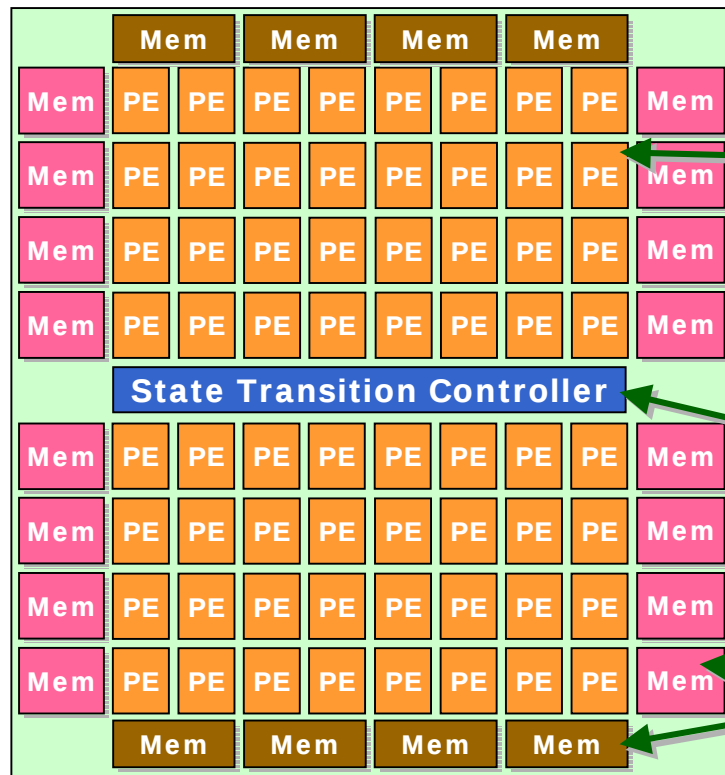
- ❖ Dynamically Reconfigurable Processor
=> DRP
- ❖ H/W-like performance by customized datapath configurations using array of PEs
- ❖ S/W-like scalability by dynamic reconfiguration of customized datapath planes
- ❖ C compiler is available
 - Based on in-house high-level synthesis tool, called *Cyber*

DRP Core (variable array size)



DRP Architecture

DRP Core (variable array size)



- ❖ Array of byte-oriented processing elements
- ❖ Fully programmable inter-PE wiring resources

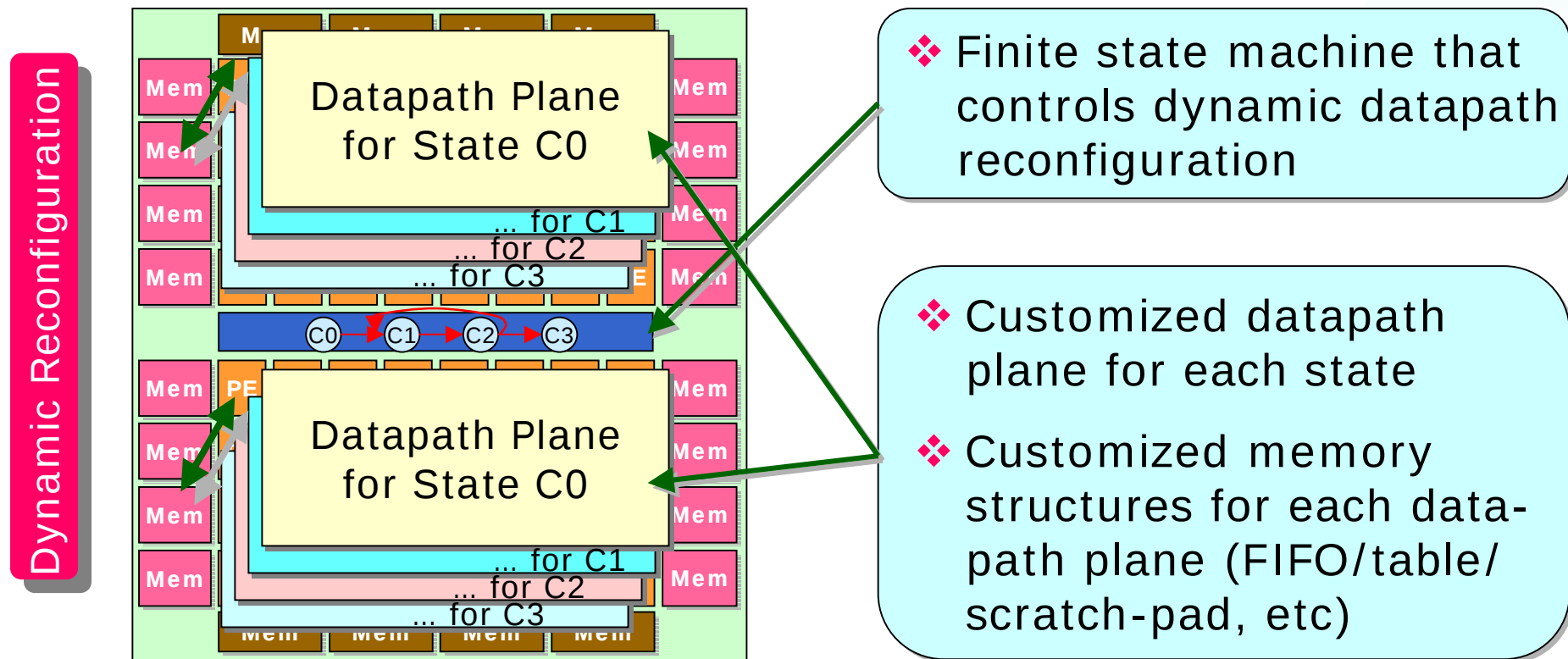
- ❖ A simple sequencer

- ❖ Array of configurable data memories

- ❖ Fine-grained processor-based programmable array architecture
- ❖ Architected for stream data processing, such as NW packet, motion/still picture, and wireless data streams, etc.

DRP Architecture

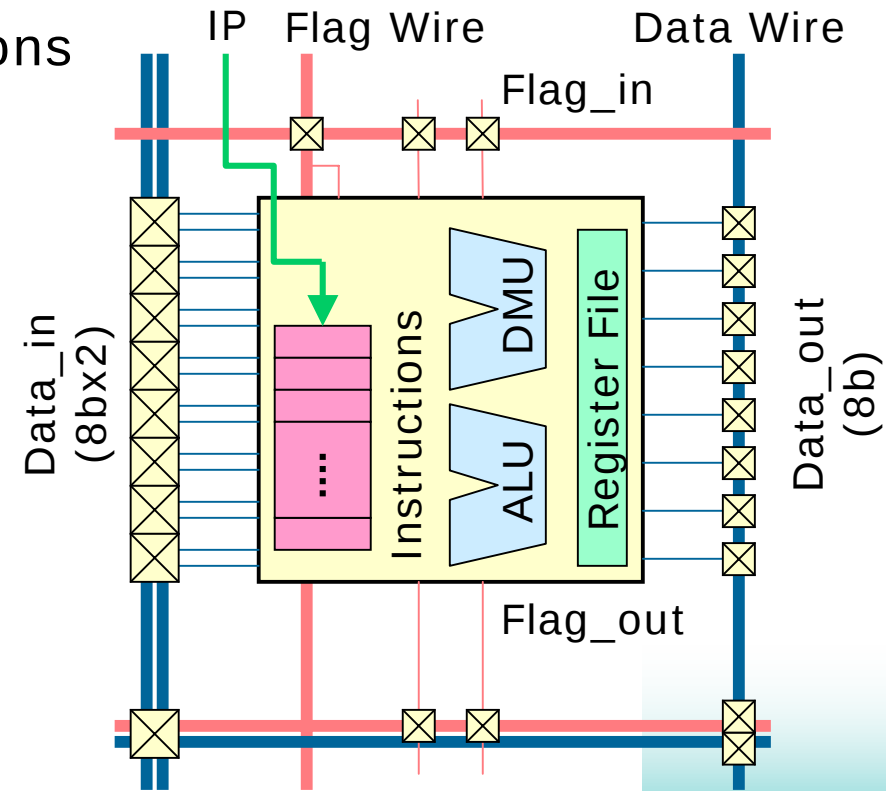
DRP Core (variable array size)



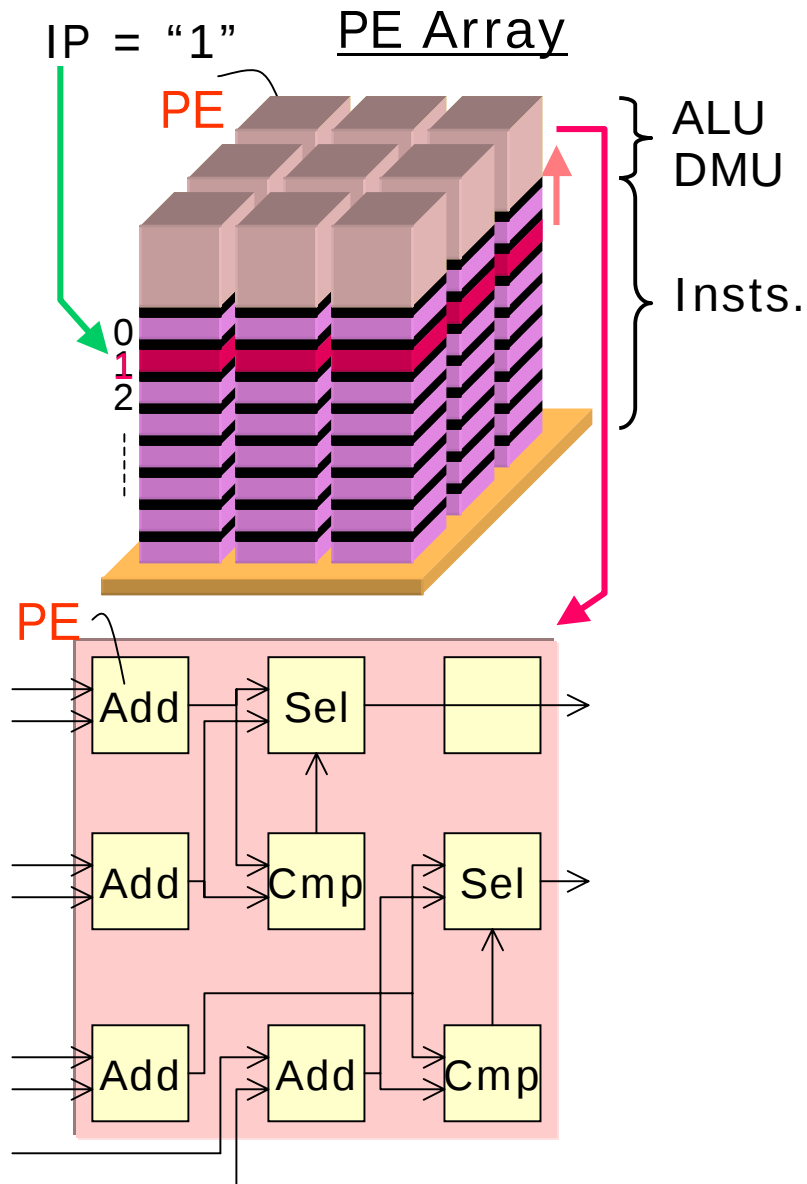
- ❖ Sequencer is built-in, and is decoupled from PE array
- ❖ Datapath is dynamically reconfigured during runtime
- ❖ Customized datapath/memory organization at any time

Processing Element (PE)

- ❖ ALU: ordinary byte arithmetic/logic operations
- ❖ DMU (data management unit): handles byte select, shift, mask, constant generation, etc., as well as bit manipulations
- ❖ An instruction dictates ALU/DMU operations and inter-PE connections
- ❖ Source/destination operands can either from/to
 - its own register file
 - other PEs (*i.e.*, flow through)
- ❖ Instruction pointer (IP) is provided from STC (state transition controller)



Computation on DRP Core



❖ Instruction Pointer(IP) from STC identifies a datapath plane

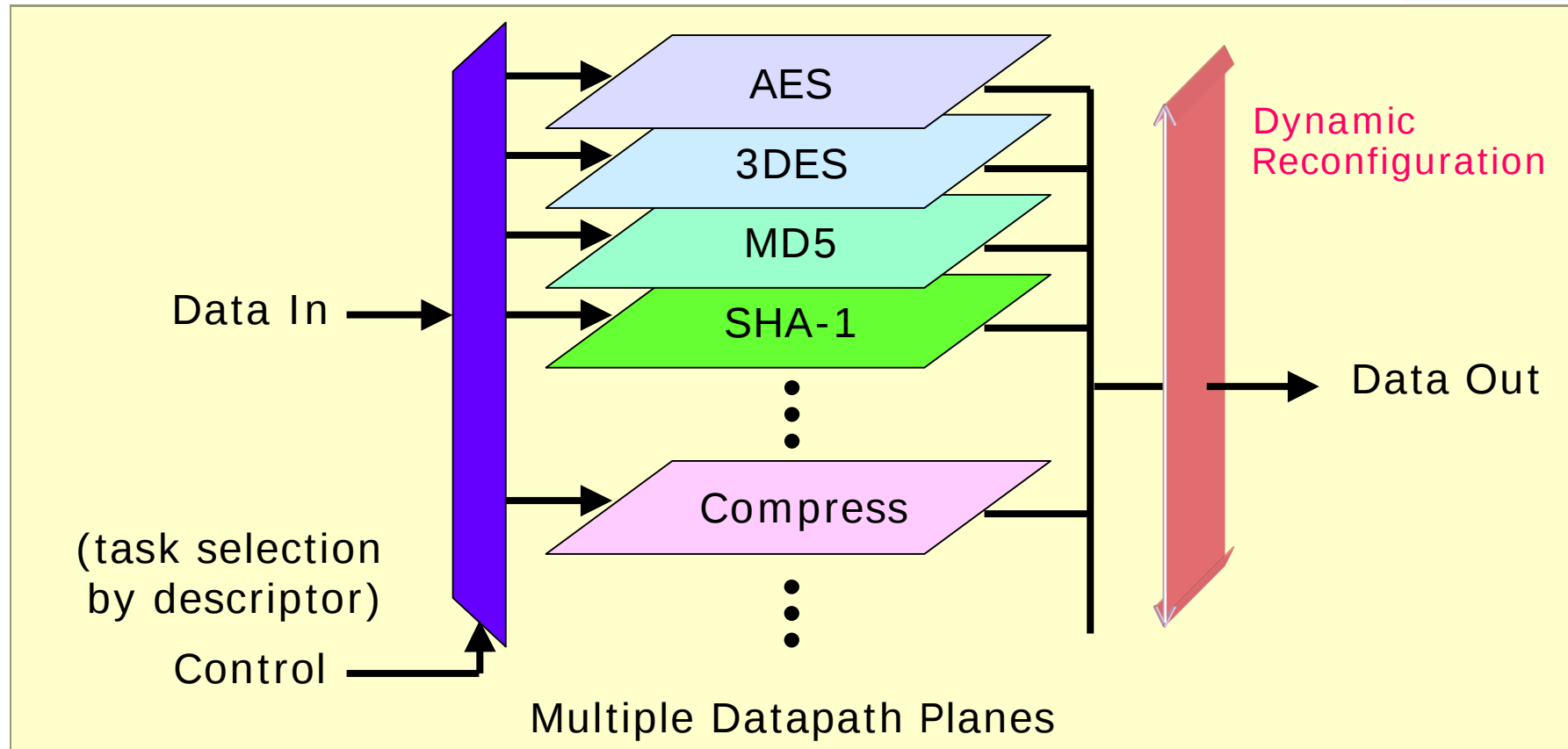
❖ Spatial computation with using a customized datapath plane

❖ When IP changes, datapath plane switches instantaneously

➤ PE instructions as a collection behave like an extreme VLIW

➤ Sequencing through instructions => Dynamic reconfiguration

Example: IPSec on DRP Core

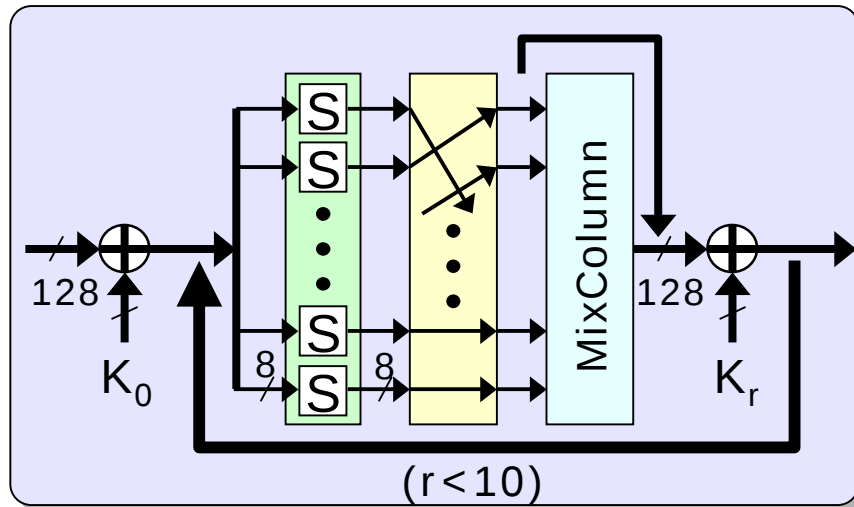


Different datapath plane for different set of algorithmic processing

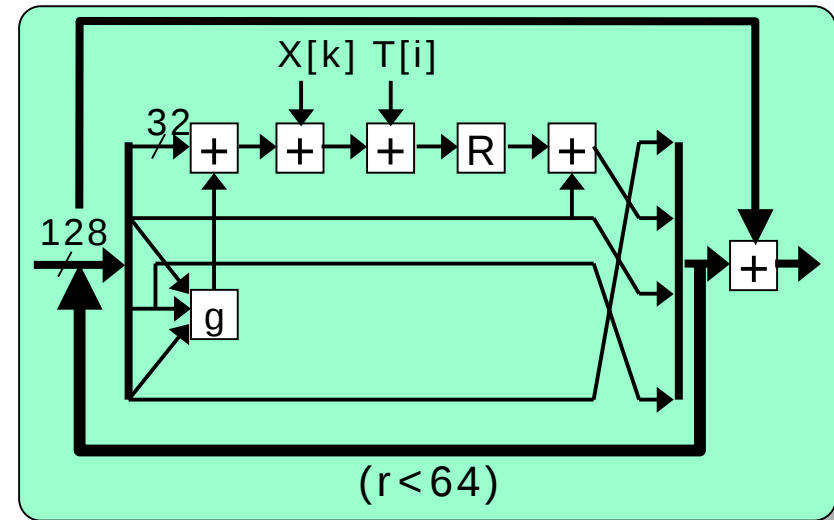
IPSEC on DRP Core - Continued

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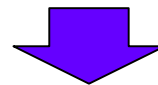
AES



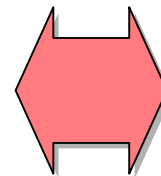
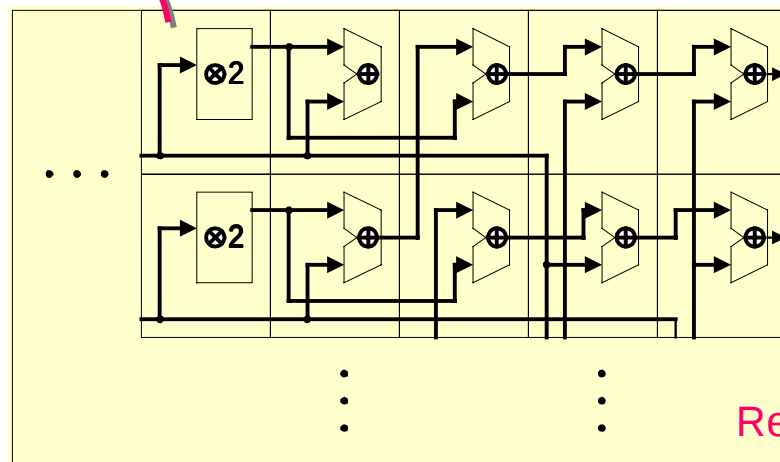
MD5



PE

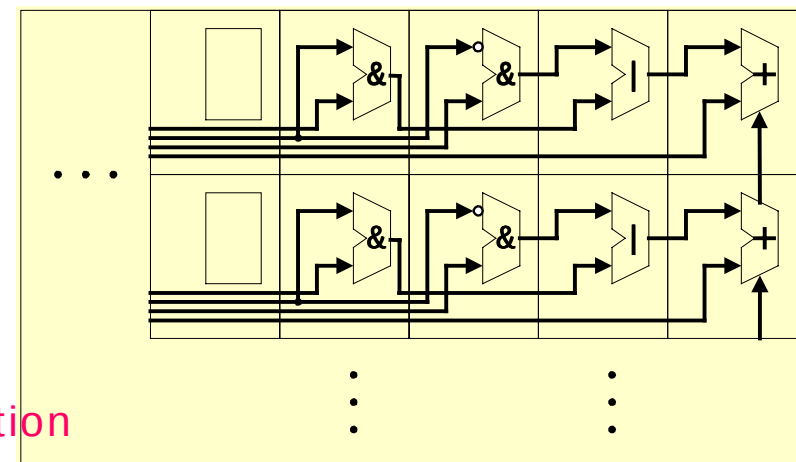
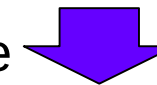


Place & Route



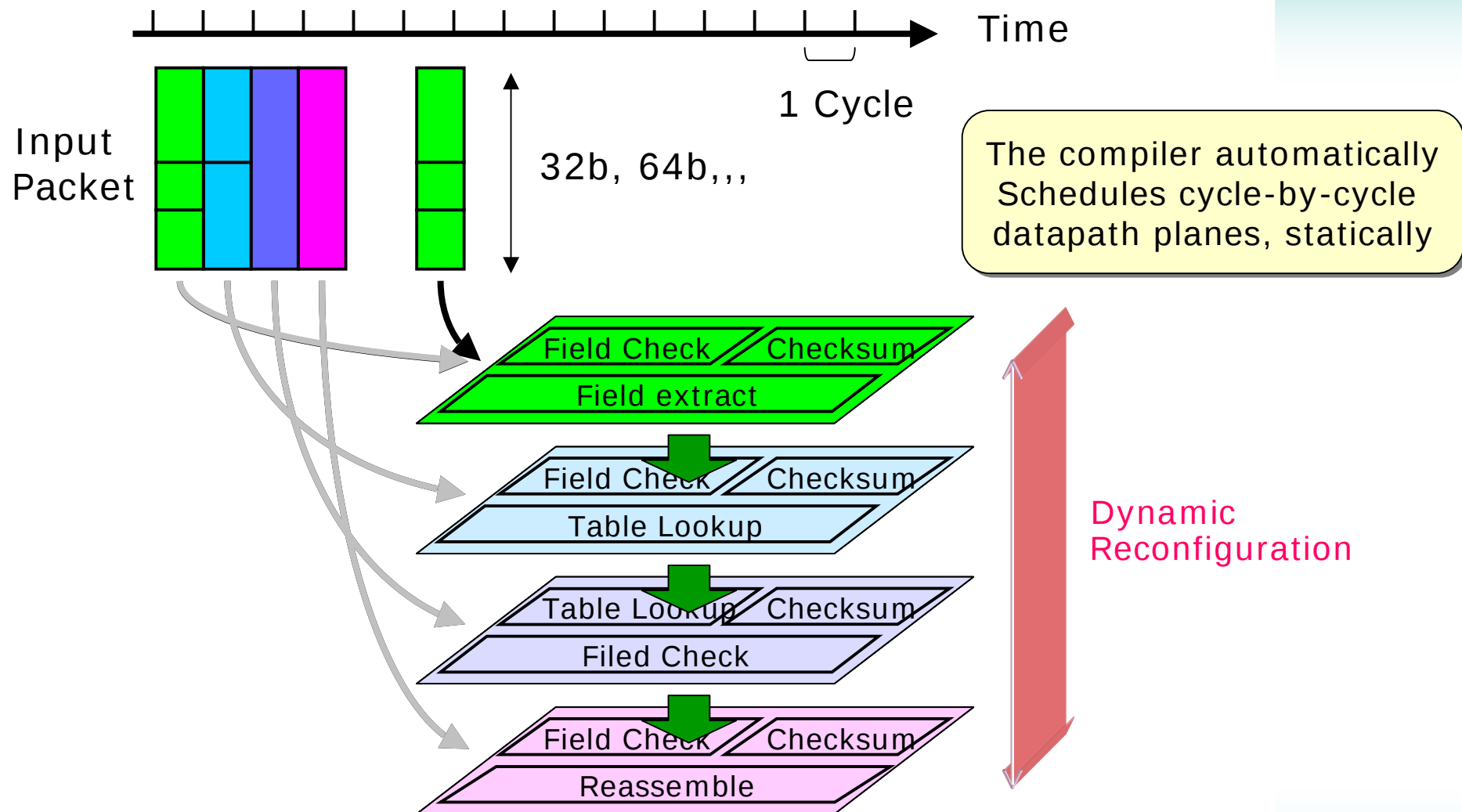
Dynamic
Reconfiguration

Place & Route



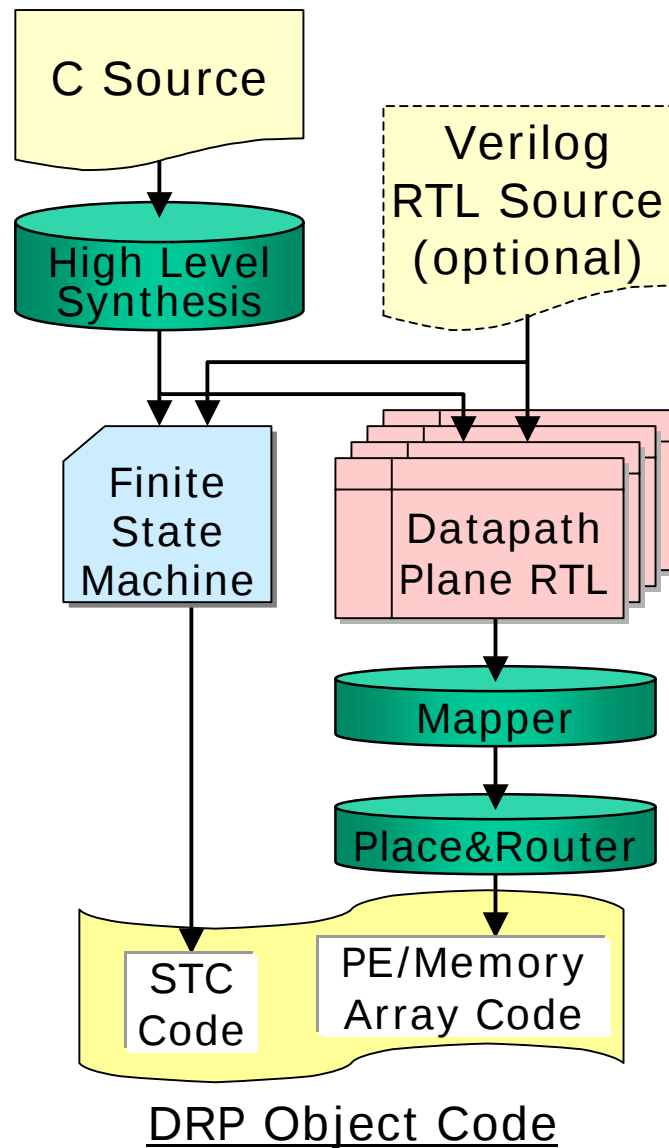
Example: Packet Processing

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Different hardware for processing respective chunk of input packet

DRP Compiler



Compile C source code
into DRP object code

- ❖ High-level synthesis: generates finite state machine and associated datapath planes from the C source code
 - Modify *Cyber* to cope with DRP architectural features
- ❖ Mapper: maps RTL for each datapath plane to individual PEs and memories
- ❖ Place&Router: physically locate the PEs and memories, and mutually connect them

Programming Model

C source code

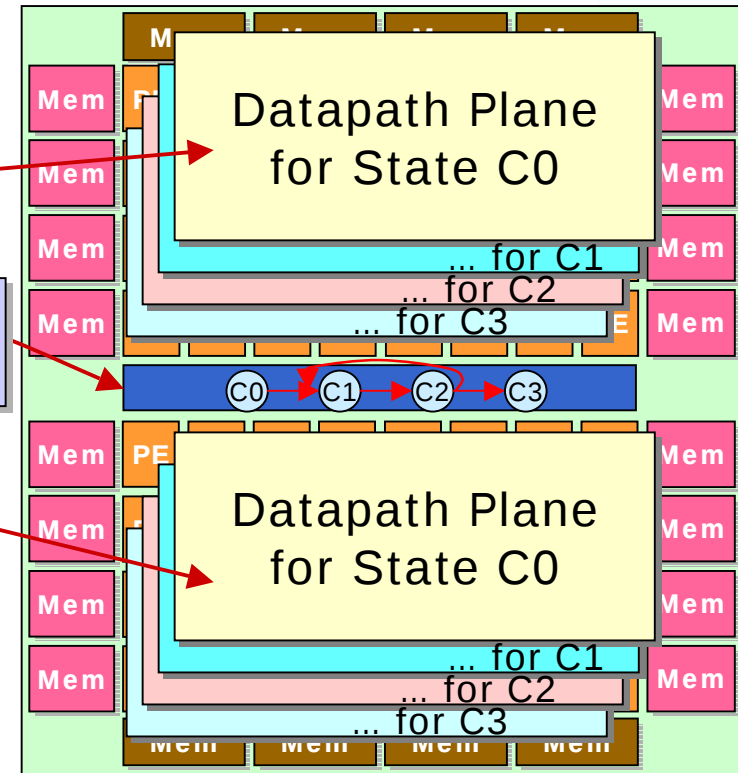
```
if ((a>b)|| (c!=d))
    DataPath_Func_1();
else if (e==f)
    DataPath_Func_2();
else
    DataPath_Func_3();
```

Bit-oriented
Processing

Control Structure
=>
Finite State Machine

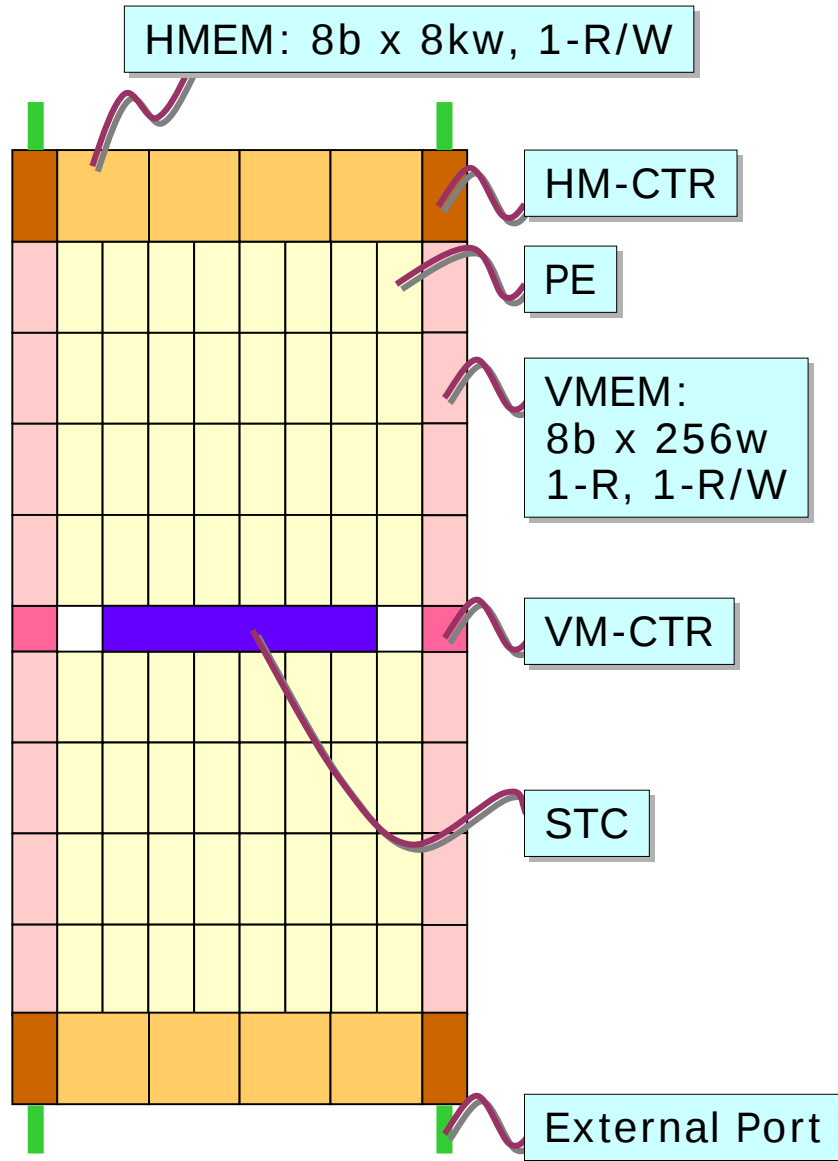
Byte-oriented
Processing

DRP-Core



- ❖ “Finite state machine with datapath planes”
- ❖ DRP is architected based on a clear picture of how C code is compiled into hardware

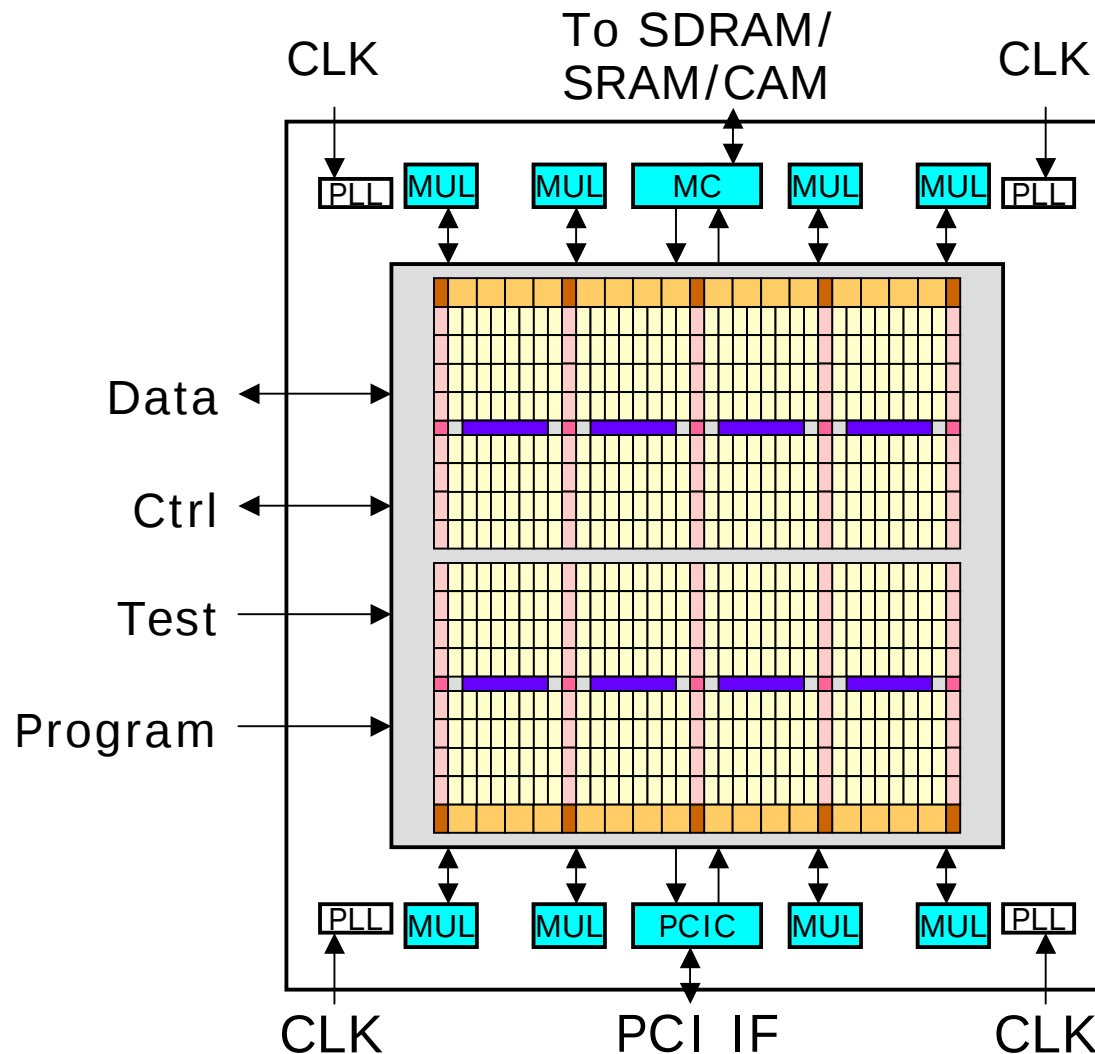
1st Implementation: DRP Tile



- ❖ Tile is a minimum unit of DRP array
 - 8x8 PEs
 - 1-p HMEM and 2-p VMEM
- ❖ A PE stores 16 instructions
- ❖ Reconfiguration of a tile is controlled by its STC
- ❖ External port can attach extended functional units, such as
 - External memory controllers
 - Peripheral bus controllers (like PCI)
 - Complex operation units (like multiplier, divider, etc)

1st Silicon: DRP-1

DRP-1 Chip Blockdiagram

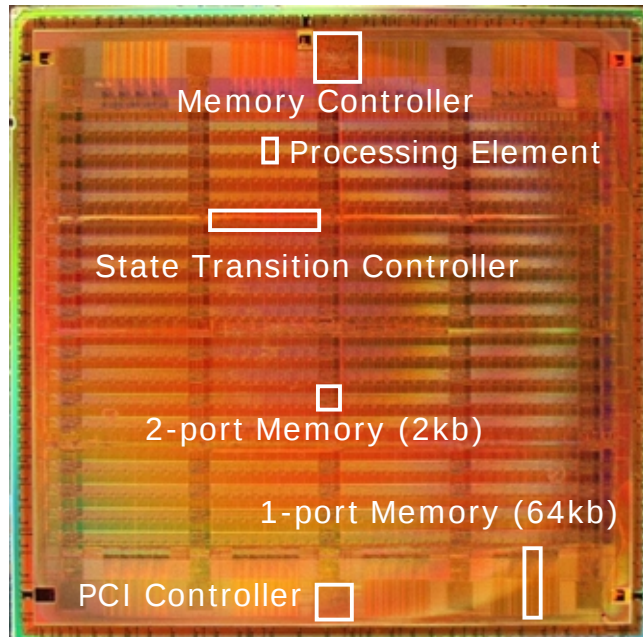


- 8 Tiles. 512 PEs, 160kb VMEMs, 2Mb HMEMs
- 8 STCs
- 8 32b Multipliers
- 1 SDRAM/SRAM/CAM controller
- 1 PCI controller
- 4 PLLs

- PE
- VMEM
- HMEM
- VM-CTR
- HM-CTR
- STC
- Extended Functional Units

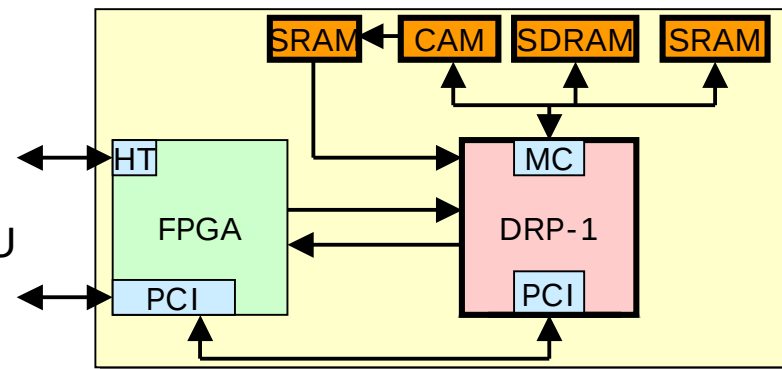
Current Status: HW Platform

DRP-1 Die Photo



- ❖ 0.15 μ m CMOS 8-AI process, 696-pin TBGA package
- ❖ 33-133MHz clock
- ❖ 22M logic transistors (44K for a PE)
- ❖ 1.5Mb configuration memories
- ❖ Fully functional w/o respin

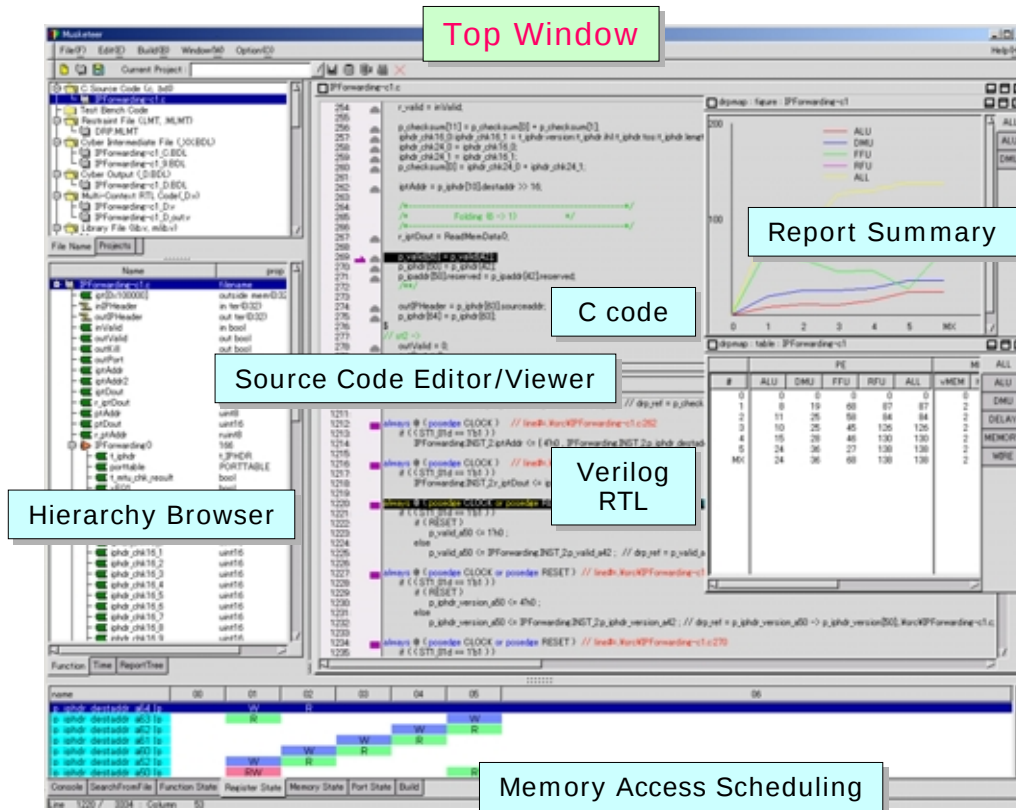
To/from
Host MPU



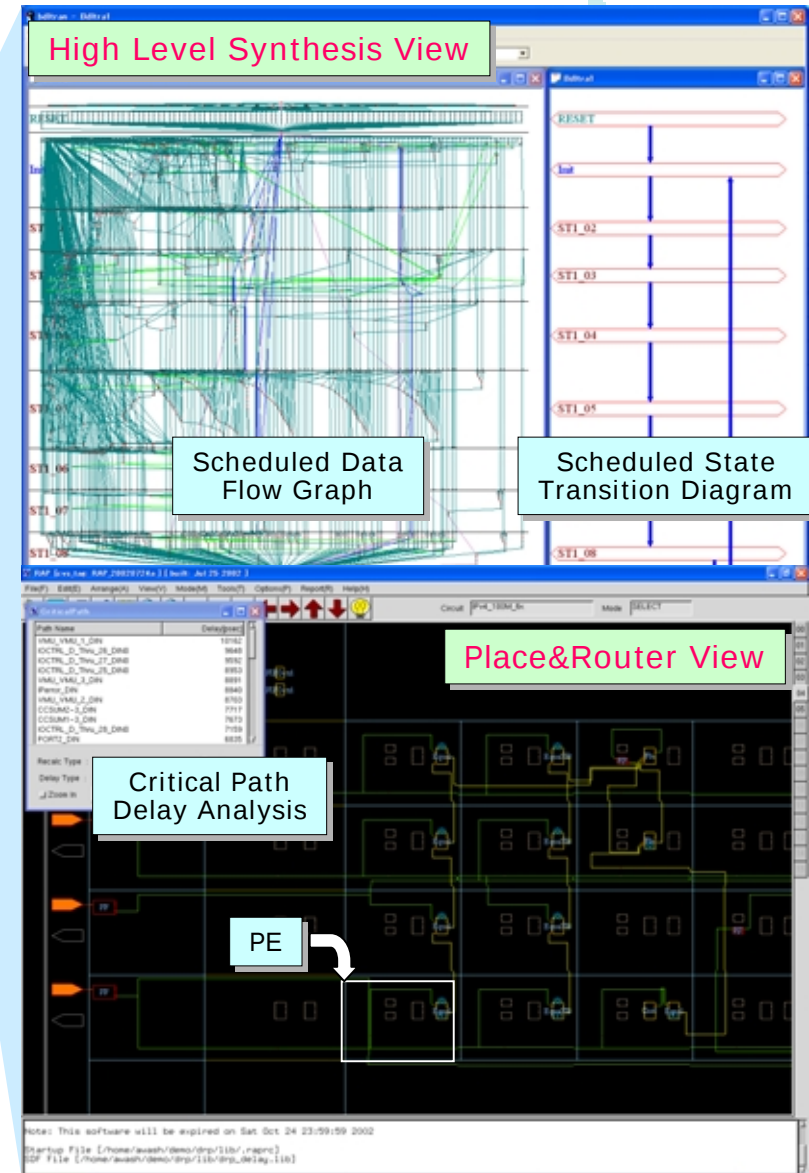
- ❖ DRP-1 board is now up and running
 - Demonstration of DRP as an off-loading engine to a host MPU through PCI or HyperTransport I/F

Current Status: Compiler Suite

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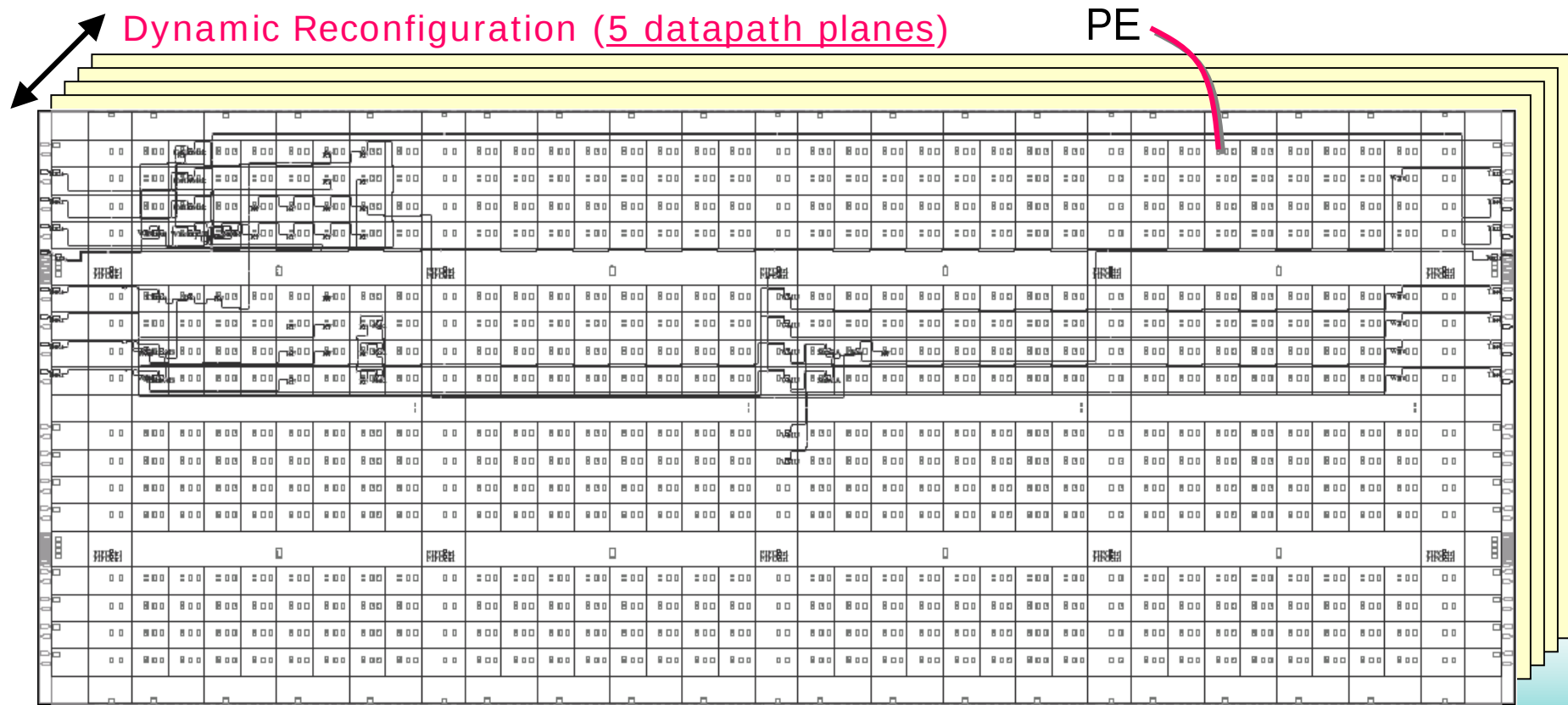


DRP application development system (HW+SW) is ready

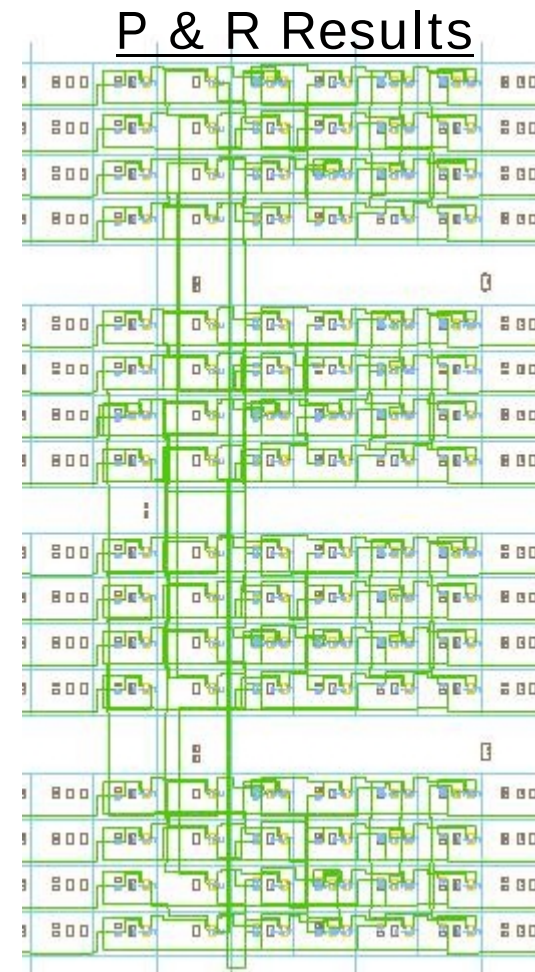
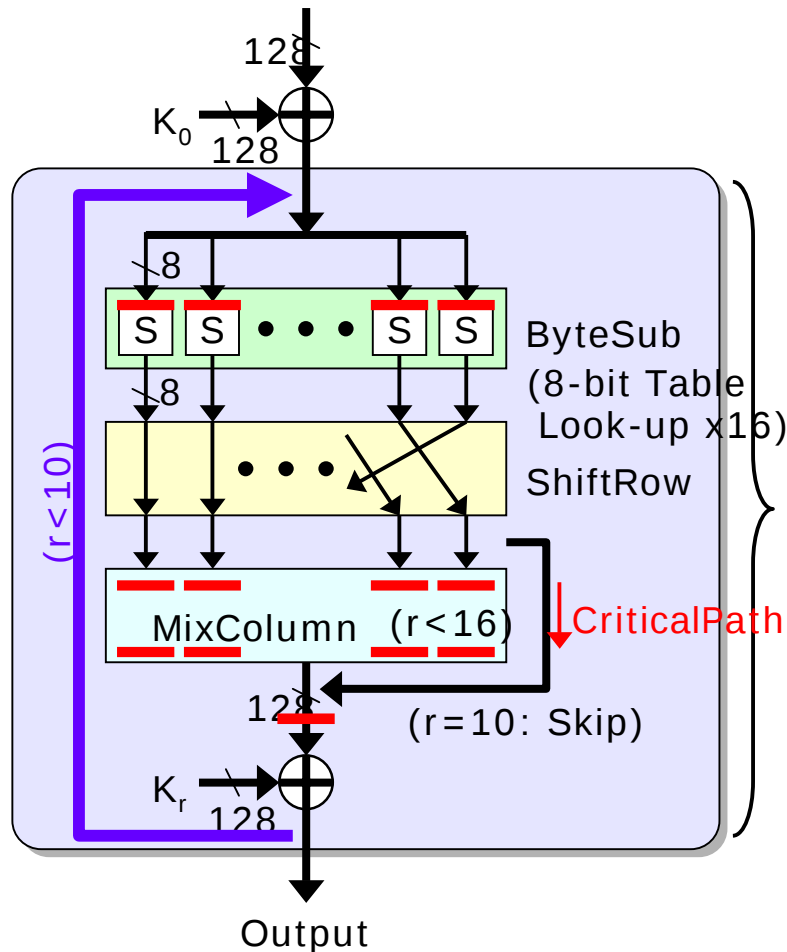


Example: IPv4 Packet Forwarding

- ❖ RFC1812 (Header check, checksum, longest prefix match, header modification)
- ❖ 8x16 PEs, 100MHz Operation
- ❖ Table look up: 5 round index search using 100MHz SRAM
- ❖ 6 pipe-stages, 5-cycle each. $100\text{MHz} \times 5\text{Cycle} \Rightarrow 20\text{M Packet/second}$



Example: AES Encryption



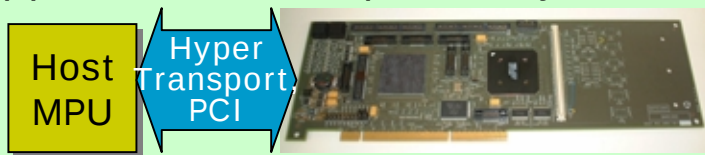
- ❖ Critical Path: 6.6ns. Throughput: 1.76Gbps (ECB mode)
- ❖ Fits in 8x16 PEs, single context (80 ALU/DMUs, 96 RFUs, 16vMEMs)

DRP Roadmap

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DRP-based Product

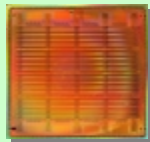
Application development system



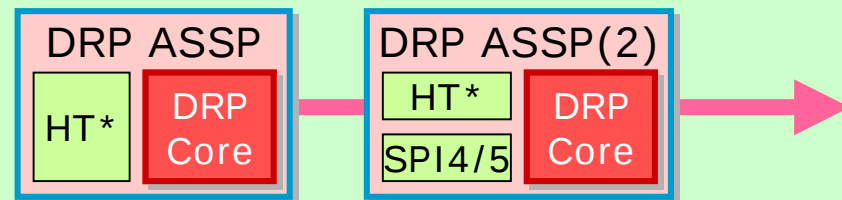
Embedded core for NEC ASICs

- ❖ Faster time to market
- ❖ Longer time in market
- ❖ Lower risk ASIC design

DRP 1st Silicon

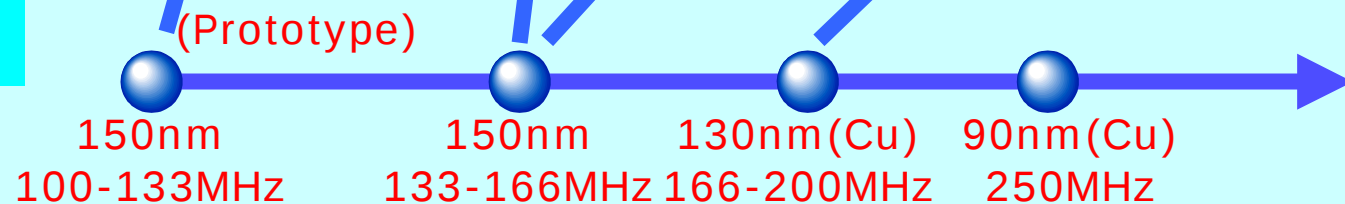


Programmable off-load engine to host MPU



* PCI-X/Express may also be available

DRP Core



Year

2002

2003

2004

Summary

- ❖ DRP architecture features
 - Array of PEs and memories for programmable datapath
 - FSM-based sequencer, a state transition controller, for the control of dynamic datapath reconfiguration
- ❖ DRP application development system (DRP-1 chip & board, and the compiler suite) is ready
- ❖ Sample application development is underway
- ❖ Launch of DRP-based product planned during Y'03